Pearson Webcast Series



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Understanding Oracle Explain Plans

Presented by **Dan Hotka**, Oracle ACE Director

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livelessons Oracle SQL Performance Tuning for Developers

Dan Hotka



About Us

Pearson

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Welcome. Our Agenda Today.

- •Explain Plan Tools
 - •Xplan.display
 - •JS Tuner
- •Understanding Explain Plans •Reading the Explain Plan •What does it mean
- Q&A
- Wrap-up & Resources

- •DBMS_XPLAN.DISPLAY
 - •Available with Oracle 8.1.7+
 - •Comes with Database
 - •Used via SQL*Plus
- •JS Tuner
 - •Available for Oracle8+
 - Download from www.DanHotka.com
 - •Executable Jar File

•PLAN_TABLE

•Oracle10g+ - use PLAN_TABLE\$

•Owned by SYS

•Have DBA make public synonym to: PLAN_TABLE

•Make sure to drop your PLAN_TABLE

•Use the 'explain plan for' syntax to populate this table

•Use tools to populate this table (TOAD, SQL Developer)

•Use SHOW_PLAN.sql to display contents

•Available on my website

Using the Tools

•All tools use the PLAN_TABLE

Oracle10g+

Automatically Created

•Migration:

•Use PLAN_TABLE\$

•Autotrace will give warning if not using correct PLAN_TABLE

Column Pk I Data Type STATEMENT ID VARCHAR2 (30) TIMESTAMP DATE REMARKS VARCHAR2 (80) OPERATION VARCHAR2 (30) OPTIONS. VARCHAR2 (255) OBJECT NODE VARCHAR2 (128) OBJECT_OWNER VARCHAR2 (30) OBJECT NAME VARCHAR2 (30) OBJECT_INSTANCE NUMBER OBJECT TYPE VARCHAR2 (30) OPTIMIZER VARCHAR2 (255) SEARCH_COLUMNS NUMBER ID NUMBER PARENT ID NUMBER POSITION NUMBER COST NUMBER CARDINALITY NUMBER BYTES NUMBER OTHER TAG VARCHAR2 (255) PARTITION START VARCHAR2 (255) PARTITION_STOP VARCHAR2 (255) PARTITION ID NUMBER OTHER LONG DISTRIBUTION VARCHAR2 (30) CPU COST NUMBER IO COST NUMBER TEMP SPACE NUMBER ACCESS_PREDICATES VARCHAR2 (4000) FILTER PREDICATES VARCHAR2 (4000)

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Expla	ined.					
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×	7	INDEX FAST FULL SCAN	A_USER0	1000	6000	3	(0)	00:00:01	

7 - filter("A"."STATUS"='OPEN')

23 rows selected.

DBMS_XPLAN

🚢 Oracle SQL*Plus

File Edit Search Options Help

1 SELECT /** GATHER_PLAN_STATISTICS */ count(*) 2 from A, B, C 3 WHERE A.STATUS = B.STATUS 4 AND A.B_ID = B.ID 5 AND B.STATUS = 'OPEN' 6 AND B.ID = C.B_ID 7* AND C.STATUS = 'OPEN' SQL> /

COUNT(*)

10000

SQL> Select plan_table_output

2 From table(dbms_xplan.display_cursor(FORMAT=>'ALLSTATS LAST'));

PLAN_TABLE_OUTPUT

SQL_ID 49563xqn62buu, child number 0

SELECT /*+ GATHER_PLAN_STATISTICS */ count(*) from A, B, C WHERE A.STATUS = B.STATUS AND A.B_ID = B.ID AND B.STATUS = 'OPEN' AND B.ID = C.B ID AND C.STATUS = 'OPEN'

Plan hash value: 2966481601

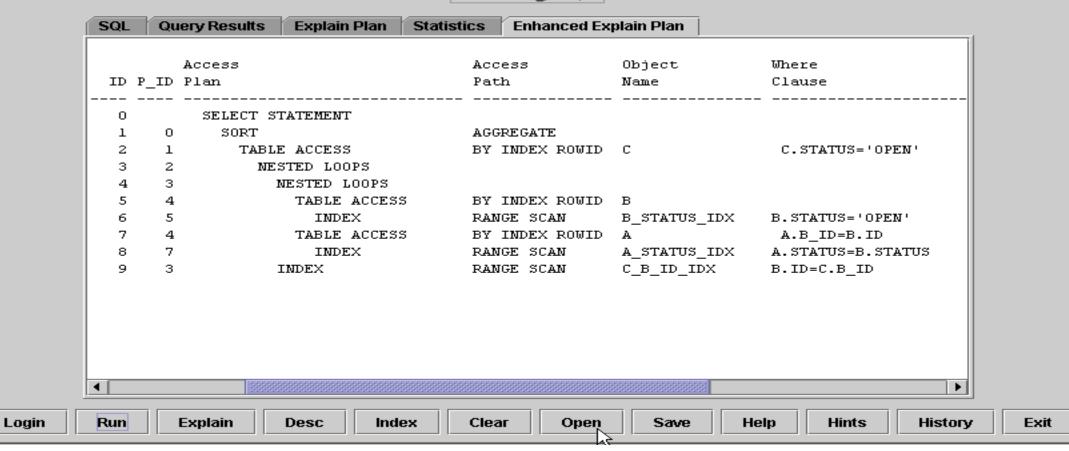
Id	Operation	Name	Starts	E-Rows	A-Rows	A-Time	Buffers OM	 em 1Mem Used-Mem
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•JSTuner

- •Incorporates SQL Scripts with enhanced Trace
- •Index Info
- Includes index statistics
- •Enhanced Explain Plan
 - •Works with V8+
 - •Works with rule-based optimizer!
- •Available from www.DanHotka.com



user0/*****@oraxp9i



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Login

Run

Explain

Desc

Index

Clear

Open

Save

Help

Hints

Exit

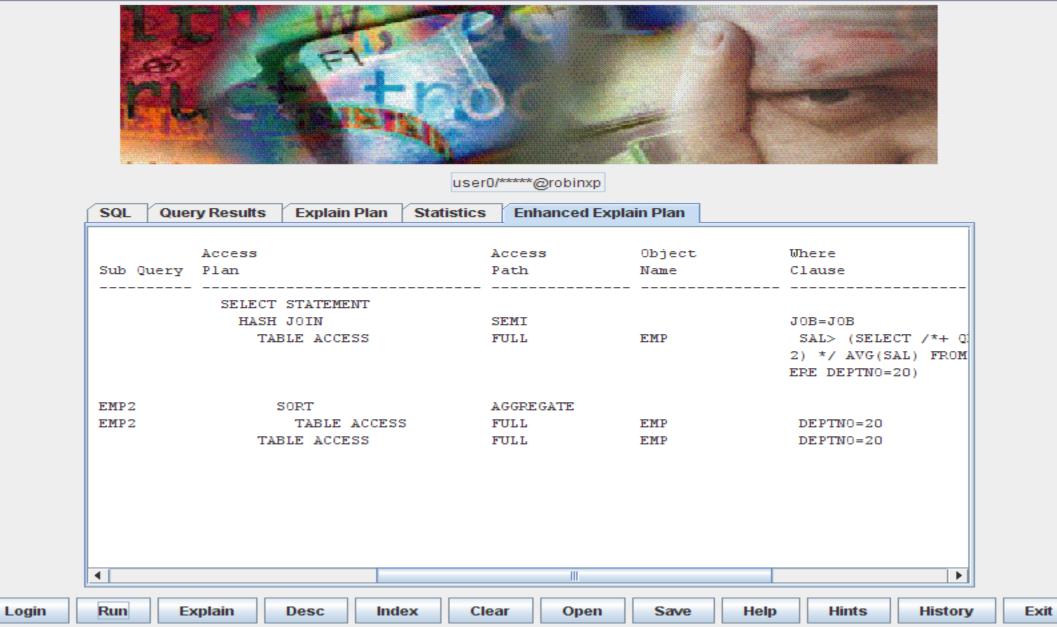
History

Hints

2 9 3 f 4 0 5 a 6 /	select From e where and sa / ined.	enar mp job i 1 > (n emp1 where deptno = 20) rom emp2 where deptno = 20)
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_	3	2	SORT	AGGREGATE		EMP2_SUBQUERY
3	4	3	TABLE ACCESS	FULL	EMP2	EMP2_SUBQUERY
3	5	1	TABLE ACCESS	FULL	EMP1	
rows	5 sele	cted.				

🕌 JS Tuner Console





Understanding Explain Plans

•Oracle:

•Hard Parses the SQL

Checks SQL syntax

Checks for available indexes/stats

•Reads from bottom to the top

Arrives at an Execution Plan

Decides how it will access the tables and indexes

•Executes the SQL

•Oracle9/10 – peeks once at bind variables

•Oracle11 – tracks explain plans with various bind vars

Called 'Adaptive Cursor Sharing'

•Runs the Execution plan

•Explain Plan

- •Visualizes the execution plan
- •Uses the Plan_Table
 - •Which varies slightly from release to release
 - •Utlxplan.sql
 - •<Oracle Home>/RDBMS/ADMIN
- Used heavily to tune SQL

•How does it work?

•Reads from bottom up

•Syntax checks/tracks available indexes

•<u>ALL</u> queries begin with a table access

•Regular queries

generally does the table joins first

Utilizes WHERE clause predicates

Partitioned queries

generally accesses the partitions first utilizing WHERE clause predicates

•Then performs the table joins

•Sometimes the tables are again

•Including remaining WHERE clause predicates

•How does it work?

•RBO - follows rules for index selection

•Arrives at an execution plan in single pass

•CBO - tries a variety of combinations of table order/where clause predicates

Combinations called `permutations'

•Generally # perms = # tables * # where predicates

•Regular queries:

Uses lowest cost method

•*** Queries in this course are regular queries unless otherwise noted

•Partitioned queries (covered in separate section)

• SQL that accesses 1 or more partitioned objects

•Uses fastest access method (elapse time)

•RBO – produces explain plans in a single pass
•Hard parsing was not the issue
•Size of the library cache was the issue

•CBO – produces multiple iterations looking for the lowest cost

•These iterations are called permutations and query transformation

•CBO – Brute Force

- •CBO is a code set
- •Gets a SQL
- •Returns Cost #'s and Execution Plans
- •Oracle passes next permutation to CBO
- •Oracle passes query rewrite to CBO

•IS overall cost (cost # at line 0) < prior permutation?

•True – tosses prior and this one becomes the one to beat

Permutations

- •CBO tries various table combinations
- •CBO costs out the 3 join types
 - Nested Loops
 - •Merge Join
 - •Hash Join
- •CBO costs out the various where clause predicates associated with each table

•Query Transformations

- •Introduced in Oracle9
 - •Complex View Merging
 - Converts views to joins
 - Subquery Unnesting
 - Converts subqueries to inline views
 - •Join Predicate Push Down
 - Moves where clause predicates into subquery

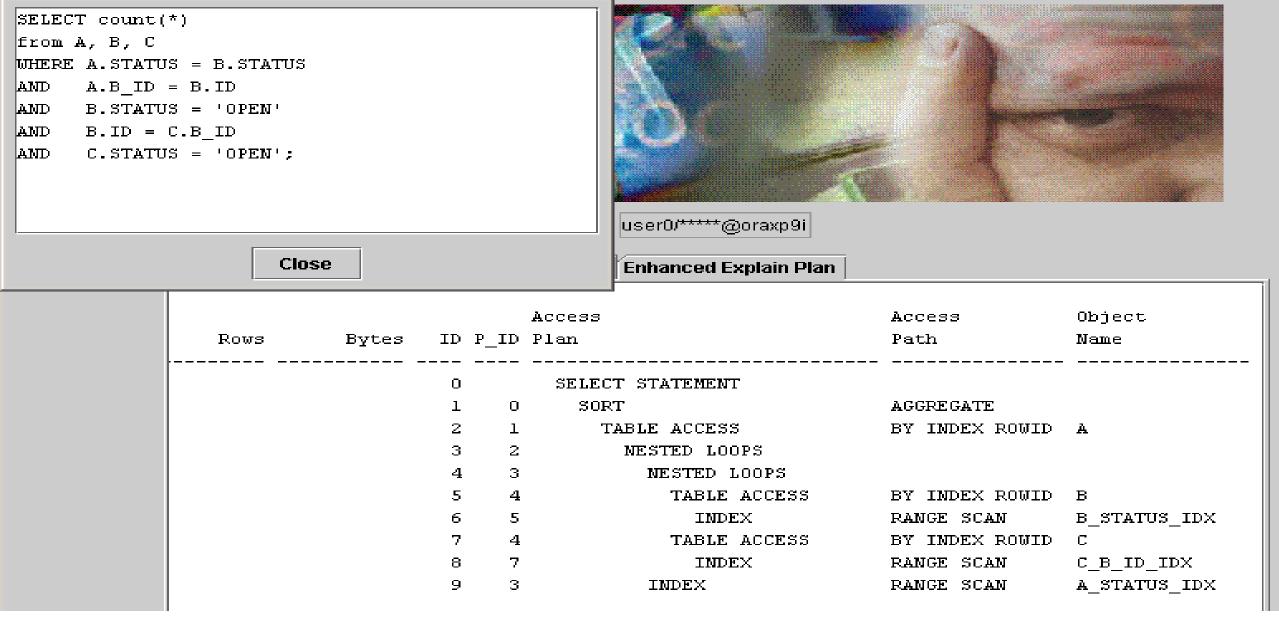
-Hints for each item above

- No_Query_Transformation prevents this behavior
- -Discussed later in this chapter

•Inner steps produce result sets

- These `intermediate' result sets are not visible
- •They are like temporary tables
- •They are passed to the next step in the execution plan
- •IF passed to a join step
 - •They become the Outer Table
- •The result set at Step 0 is the final result set
 - •This result set is then passed to the cursor area

Current SQL



🌺 Current SQL SELECT count(*) from A, B, C WHERE A.STATUS = B.STATUS $A.B_{ID} = B.ID$ AND. AND B.STATUS = 'OPEN' $B.ID = C.B_ID$ AND AND C.STATUS = 'OPEN'; user0/*****@oraxp9i Close Enhanced Explain Plan Where Access Access Object ID P ID Plan Path Name Clause SELECT STATEMENT 0 SORT 0 AGGREGATE 1 BY INDEX ROWID A 2 l TABLE ACCESS A.B ID=B.ID з. 2 NESTED LOOPS 4 3 NESTED LOOPS 4 TABLE ACCESS BY INDEX ROWID B 5 6 5 INDEX RANGE SCAN B_STATUS_IDX B.STATUS='OPEN'

BY INDEX ROWID C

RANGE SCAN

RANGE SCAN

C.STATUS='OPEN'

B.ID=C.B_ID

A.STATUS=B.STATUS

C_B_ID_IDX

A_STATUS_IDX

TABLE ACCESS

INDEX

INDEX

4

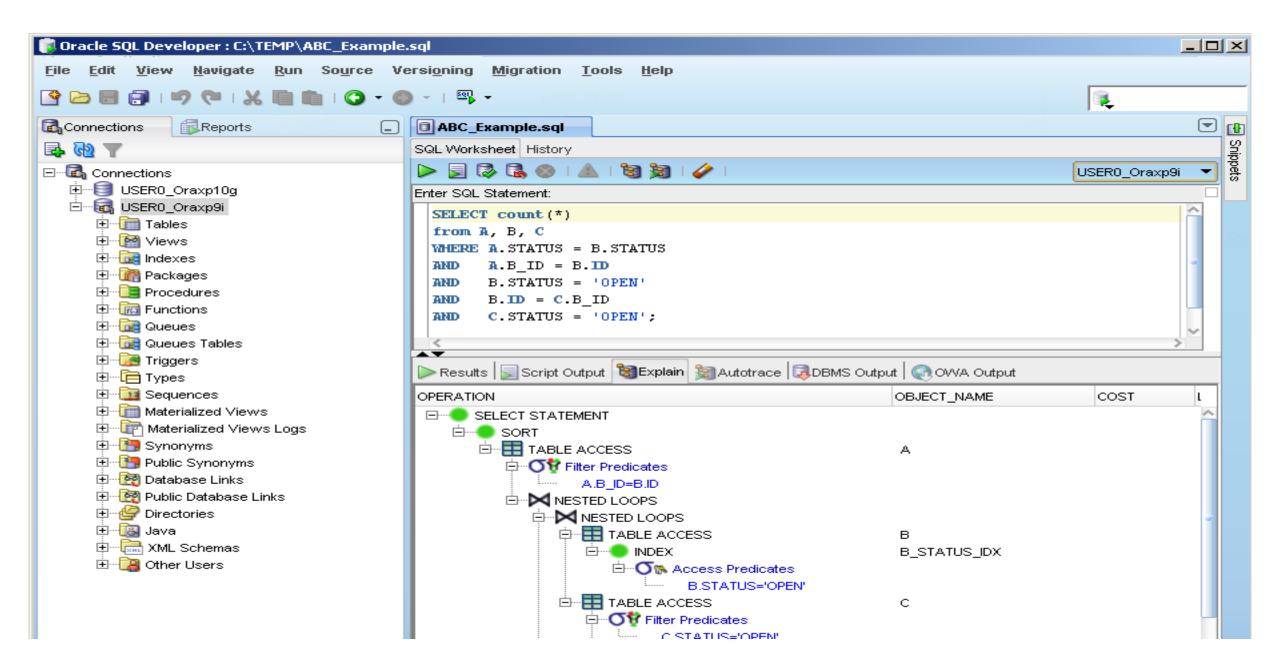
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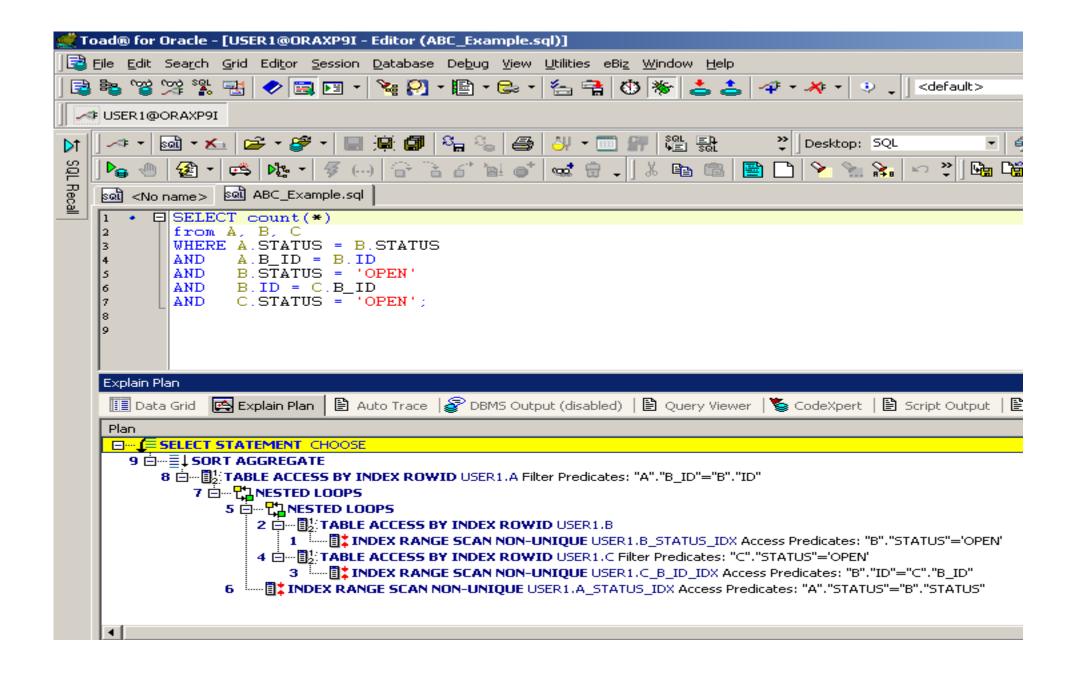
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9





- •<u>Access Rule Description</u>
- •AND-EQUAL Index values will be used to join rows.
- •Access Predicate works with ROWID's, tyically from indexes.
- •Filter Predicate FILTERs apply 'other criteria' in the query to further qualify the matching rows. The 'other criteria' include correlated sub queries, and HAVING clause.
- •TABLE ACCESS When not associated with a join condition, they act like Filter...processing additional Where Clause predicates.
- •VIEW OF Processed SQL from a view. **IF on a join condition, Oracle converted it to a view/subquery during query transformation
- •INTERNAL_FUNCTION This typically means that you have a data type mismatch.

- •Access Rule Description
- •CONCATENATOIN statement has a union clause

•INDEX (UNIQUE) SQL statement utilized a unique index to search for a specific value.

•INDEX (RANGE SCAN) SQL statement contains a non-equality or BETWEEN condition.

•INLIST ITERATOR SQL statement has an IN clause, or, values being treated as an IN clause

•TABLE ACCESS (FULL) All rows are retrieved from the table without using an index.

•TABLE ACCESS (BY ROWID) A row was retrieved from a table based on the ROWID of the row.

•Access Rule Description

•HASH JOIN SQL statement initiated a hash-join operation.

•MERGE JOIN SQL statement references two or more tables, sorting the two result sets being joined over the join columns and then merging the results via the join columns.

•MERGE JOIN (CARTESIAN) SQL statement references two or more tables but without a joining column (generally not a good thing)

•NESTED LOOPS This operation is one form of joining tables, as opposed to a merge join. One row is retrieved from the row source identified by the first child operation, and then joined to all matching rows in the other table, identified in the second child operation.

•NONUNIQUE INDEX (RANGE SCAN) The RANGE SCAN option indicates that ORACLE expects to return multiple matches (ROWIDs) from the index search

•Access Rule Description

- •BITMAP CONVERSION Bitmap Index being merged
- •BITMAP MERGE Generally used with Bitmap Range Scan
- •BITMAP MINUS Bitmap Index handling a not = condition.
- •BITMAP INDEX SINGLE VALUE Bitmap index being used for single value lookup.
- •BITMAP INDEX (RANGE SCAN) Bitmap index being used for multiple value lookup.

•Access Rule Description

•PARTITIONING covered in another ppt...ask for it next time!

•SORT (ORDER BY) SQL statement contains an ORDER BY SORT (AGGREGATE) SQL statement initiated a sort to resolve a MIN or MAX type function.

•SORT (GROUP BY) SQL statement contains a GROUP BY

•Index Scans

•Unique Scan

•via root -> branch -> leaf for a single row acces:

•Range Scan

 via root -> branch -> leaf for first row, then lea to leaf for remaining rows

•Full Scan (Index-Full)

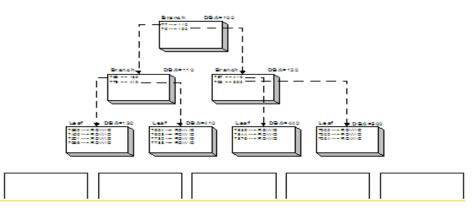
•Scans leaf blocks using single-block access

Fast Full Scan (Index-FFS)

•scans all index leaf blocks using multi-block read

•Index Skip Scan (Index-SS)

•Useful for multi-column indexes, accessing only 2nd column and first column has low cardinality



- •Oracle only joins 2 tables at a time
- •The smaller the initial result sets, the faster the whole query runs
 - •Drive off the item that will eliminate the most rows first

Nested Loops

- •Inner table looped for each row returned in outer table
- •Lg table should be outer
- •Sm table (or unique indexed lookup) should be inner
- •Rows returned to the result set that qualify the driving WHERE clause
- •Cost = outer access + (inner table access * outer cardinality)

•Merge Scan Join

- •Both tables are sorted
- •Rows are inserted into result set based on key value
- •THEN WHERE clause applied
- •Cost = outer access + inner access + sort costs

•Hash Join

- •Hashes join keys and caches object into a hash table
- •Driving table should be smaller of the 2
- •Cost = inner cost + (outer cost * inner cardinality/hash partitions)

•Nested Loop Join

•driving table

Default order(rule)

•Merge Scan Join

•sort & match

•Hash Join

•Full scans with no sorts

•Join column to row address

■1 - NESTED LOOPS | ●■1 - TABLE ACCESS [FULL] of ECO.CONTACTS | ●■2 - TABLE ACCESS [BY ROWID] of ECO.COMPANIES | ↓■1 - UNIQUE INDEX [UNIQUE SCAN] of ECO.PK_COMP_KEY[COMP_KEY]

🚍 1 - MERGE JOIN		
La1 - SORT [JOIN] La1 - TABLE ACCESS [FULL] 2 - SORT [JOIN]		
Le 1 - TABLE ACCESS [FULL]	of	ECO.CONTACTS
L=2 - SORT [JOIN]		
Lal - TABLE ACCESS [FULL]	of	ECO.COMPANIES



- •Nested Loop Join
 - •If join condition is 'ANDed', make a compound index on the inner table
 - •Inner and Outer join column should have same data type
 - •Outer Table: Larger of Result set
 - •Inner Table: Smaller of Result set
 - •Foreign key indexes helps CBO choose between nested loops and hash joins

•Merge Scan Join

large portion of rows are being joined

•Both Tables have larger result sets

•Helpful if using indexes on merged columns (rows returned in proper order)

•Hash Join

•Smaller tables being joined on `=` condition

•Outer Table: Smaller of Result Set

•Inner Table: Larger of Result Set

•Use Trace Event 10104 (Hash Area Trace) or 10053 (CBO Trace) to size correctly

•HASH_AREA_SIZE

Additional Resources

Dan Hotka

- Dan@DanHotka.com
- Website: www.DanHotka.com
- Twitter: @DanHotka



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